

Monte-Carlo-Type Techniques for Processing Interval Uncertainty, and Their Engineering Applications

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Uncertainty is Important

- In engineering, decisions are made under *uncertainty*.
- *Main source of uncertainty*: measurement errors.
- *Additional source of uncertainty*: we do not know how exactly the devices will be used.
- *Example*:
 - we have limits L_i on the loads l_i in different rooms i ;
 - we do not know how exactly these loads will be distributed; and
 - we want to make sure that our design is safe for all possible $l_i \leq L_i$.

Traditional Statistical Approach

- Traditionally, *statistical* methods are used.
- Usually, we can safely *linearize* the dependence of the desired quantities y (e.g., stress at different structural points) on the uncertain parameters x_i .
- Thus, we enable *sensitivity analysis*.
- *Problem:* for n parameters, we need n calls to the model.
- Often, the number n of uncertain parameters is huge.
- *Example:* in ultrasonic testing, we record (= measure) signal values at thousands moments of time.
- *Solution:* Monte-Carlo simulations.
- *Advantage:* the number of calls to a model depends only on the desired accuracy ε and not on n .
- So, for large n , these methods are much faster.

Formulation of the Problem

- *Problem:* in real life, we often do not know the exact probability distribution of measurement errors and/or of user loads.
- *Interval uncertainty:* often, all we know is the *intervals* of possible values of the corresponding parameters.
- *Example:* we know that the load l_i is in $[0, L_i]$.
- *Sensitivity analysis:* we can use sensitivity analysis, we can use interval techniques.
- *Problem:* for large n , this takes too long.
- *What we are planning to do:* describe Monte-Carlo type techniques for interval uncertainty.
- *Advantage:* these techniques lead to faster computations.

Formulation of the Problem

(cont-d)

- *We know:*
 - the algorithm $f(x_1, \dots, x_n)$;
 - the measured values $\tilde{x}_1, \dots, \tilde{x}_n$; and
 - the information about the uncertainty

$$\Delta x_i \stackrel{\text{def}}{=} \tilde{x}_i - x_i$$

of each direct measurement.

- *We must estimate:* the uncertainty $\Delta y = \tilde{y} - y$ of the algorithm's output.

Types of Uncertainty

- *Idea.* We must know:
 - what are the possible values of Δx , and
 - how often can different possible values occur.
- *Ideal case – full info:* we know the cdf $F_i(t)$ for each variable x_i (and we know that x_i are independent).
- *Interval case – no info about probabilities:* we know the interval $[\underline{x}_i, \bar{x}_i]$ of possible values of each x_i .
- *General case: partial info:* we know the intervals $[\underline{F}_i(t), \bar{F}_i(t)]$ that contain $F_i(t)$ (*p-boxes*).
- *Important case – DS:* we know that $x_i \in [\underline{x}_i^{(k)}(t), \bar{x}_i^{(k)}(t)]$ with probability $p_i^{(k)}$.
- *Comment:* we may have different info for different x_i .
- *Comment:* we may have dependent x_i .

Black Box

- *Traditional approach:* apply f step-by-step to the corresponding “uncertain numbers”.
- E.g.: probability distributions, intervals, p-boxes.
- *Problem:* in several practical situations, f is given as a *black box*:
 - we do not know the sequence of steps forming f ;
 - we can only plug in different values into f and see the results.
- *Examples:*
 - *commercial software:* safeguard vs. competitors;
 - *classified security-related software:* safeguard vs. adversary.
- *Additional problem:* sometimes, f takes so much time that it is only possible to run it a few times.

Sensitivity Analysis: Reminder

- *When applicable:* $f(x_1, \dots, x_n)$ is monotonic (increasing or decreasing) with respect of each of its variables.
- *Example:* linearizable f .
- *Algorithm:*
 - Compute $\tilde{y} = f(\tilde{x}_1, \dots, \tilde{x}_n)$.
 - For each i , compute

$$y'_i = f(\tilde{x}_1, \dots, \tilde{x}_i, \tilde{x}_i + h, \tilde{x}_{i+1}, \dots, \tilde{x}_n) \quad (h > 0).$$
 - Compute $\underline{y} = f(x_1^-, \dots, x_n^-)$ and $\bar{y} = f(x_1^+, \dots, x_n^+)$, where:
 - * if $y'_i \geq \tilde{y}$, then $x_i^- = \underline{x}_i$ and $x_i^+ = \bar{x}_i$;
 - * if $y'_i < \tilde{y}$, then $x_i^- = \bar{x}_i$ and $x_i^+ = \underline{x}_i$.
- *Problem:* we need $n + 3$ calls to f .
- For large n and for complex f , this is too slow.

Cauchy Deviates Method

- *When applicable:* linearizable f .
- *In this case:* $[\underline{y}, \bar{y}] = [\tilde{y} - \Delta, \tilde{y} + \Delta]$, where

$$\Delta = \sum_{i=1}^n |c_i| \cdot \Delta_i, \text{ and } c_i = \frac{\partial f}{\partial x_i}.$$
- *What is Cauchy:* $\rho(x) = \frac{\Delta}{\pi \cdot (x^2 + \Delta^2)}$.
- *Why Cauchy:* if ξ_1, \dots, ξ_n are independent Cauchy w/ Δ_i , then $\sum_{i=1}^n c_i \cdot \xi_i$ is Cauchy with desired Δ .

- *Algorithm:*

- compute $\delta x_i^{(k)} = \Delta_i \cdot \tan(\pi \cdot (r_i - 0.5))$, $r_i = U[0, 1]$.
- compute $\delta y^{(k)} \stackrel{\text{def}}{=} f(\tilde{x}_1 + \delta x_1^{(k)}, \dots, \tilde{x}_n + \delta x_n^{(k)}) - \tilde{y}$;
- find Δ from the MLM:

$$\frac{1}{1 + \left(\frac{\delta y^{(1)}}{\Delta}\right)^2} + \dots + \frac{1}{1 + \left(\frac{\delta y^{(N)}}{\Delta}\right)^2} = \frac{N}{2}.$$

- *Advantage:* after $N = 200$ runs, we get 20% accuracy $0.2 \cdot \Delta$ with 95% certainty (corr. to $2\sigma_e$).

Applications: Brief Overview

- *Environmental and power engineering:*
safety analysis of complex systems.
- *Civil engineering:* building safety (f is FEM).
- *Petroleum and geotechnical engineering:* f solves inverse problem (x_i – traveltimes).
- *Results:*
 - In the environmental and civil engineering, same results as sensitivity analysis, but faster.
 - In geotechnical engineering, the dependence of the accuracy on the location and depth fits much better with the geophysicists' understanding than statistical estimates.

Limitations of Cauchy Deviate Techniques

- Cauchy deviate technique is based on the following *assumptions*:
 - that the measurement errors are *small*, so we can safely linearize the problem;
 - that we only have *interval* information about the uncertainty, and
 - that we can actually call the program f 200 times.
- *Problem*: in real-life engineering problems, these assumptions are often not satisfied.
- *What we plan to do*: we describe how we can modify the Cauchy techniques to overcome these limitations.

What If We Cannot Perform Many Iterations

- *Problem:* in many real-life engineering problems, we cannot run f 200 times.
- *Idea:* use Cauchy estimates with the available amount of $N \ll 200$ iterations, but use new formulas for Δ .
- *Fact:* for reasonable large N , $\tilde{\Delta} - \Delta$ is \approx Gaussian.
- *Solution:* $\Delta \leq \tilde{\Delta} \cdot \left(1 + k_0 \cdot \sqrt{\frac{2}{N}}\right)$ (where $k_0 = 2$) w/certainty 95%.
- *Comment:* to get 99.9% certainty, take $k_0 = 3$.
- *Example:* for $N = 50$, $\Delta \leq 1.4 \cdot \tilde{\Delta}$ – not bad.
- *Problem:* for smaller N , $\tilde{\Delta} - \Delta$ is not Gaussian.
- *Solution:* we empirically find the factor.

Dempster-Shafer Knowledge Bases:

An Idea

- *Problem:* for each i , instead of a single interval \mathbf{x}_i , we have several intervals $\mathbf{x}_i^{(k)}$ with probabilities $p_i^{(k)}$.
- *Difficulty:* even if we have 2 intervals for $n = 50$ inputs, we have an astronomical number of $2^{50} \approx 10^{15}$ output intervals.
- *Fact:* when $\mathbf{x}_i = [x_i^{\text{mid}} - \Delta_i, x_i^{\text{mid}} + \Delta_i]$, then $\mathbf{y} = [y^{\text{mid}} - \Delta, y^{\text{mid}} + \Delta]$, where:

$$y^{\text{mid}} = \bar{y} + \sum_{i=1}^n c_i \cdot (y_i^{\text{mid}} - \bar{y}_i); \quad \Delta = \sum_{i=1}^n |c_i| \cdot \Delta_i.$$
- *DS case:* we have different pairs $(y_i^{\text{mid}}, \Delta_i)$ with different probabilities.
- *Idea:* due to Central Limit Theorem, (y^{mid}, Δ) is approximately normally distributed.
- *Comment:* not *exactly* normal since $\Delta \geq 0$.

Analysis of the Problem

- *Previously:* Cauchy distribution with given Δ .
- *Characteristic function:*

$$E[\exp(i \cdot \omega \cdot \xi)] = \exp(-|\omega| \cdot \Delta).$$

- *Now:* Gaussian mixture of several Cauchy distributions, with given Δ .
- *Characteristic function:*

$$E[\exp(i \cdot \omega \cdot \xi)] = \int \frac{1}{\sqrt{2 \cdot \pi} \cdot \sigma} \cdot \exp\left(-\frac{\Delta - \mu}{2\sigma^2}\right) \cdot \exp(-|\omega| \cdot \Delta) d\Delta.$$

- *Simplified expression:*

$$E[\exp(i \cdot \omega \cdot \xi)] = \exp\left(\frac{1}{2} \cdot \sigma^2 \cdot \omega^2 - \mu \cdot |\omega|\right).$$

Algorithm

- For different real values $\omega_1, \dots, \omega_k > 0$, compute $l(\omega_k) \stackrel{\text{def}}{=} -\ln(c(\omega_k))$, where

$$c(\omega_k) \stackrel{\text{def}}{=} \frac{1}{N} \cdot \sum_{k=1}^N \cos(\omega \cdot y^{(k)}).$$

- Use the Least Squares Method to find the values μ and σ for which

$$\mu \cdot \omega_k - \frac{1}{2}\sigma^2 \cdot \omega_k^2 \approx l(\omega_k).$$

- The resulting value μ is the average Δ .
- We repeat the above algorithm twice:
 - for samples for which $y^{\text{mid}} \leq E[y^{\text{mid}}]$, and
 - for samples for which $y^{\text{mid}} > E[y^{\text{mid}}]$.
- Based on two μ 's, we compute $E[\Delta]$ and $\sigma[\Delta]$.

What About p-Boxes?

- *Known fact:* a p-box can be described as a DS knowledge base.
- *Specifics:* a p-box $[\underline{F}(t), \overline{F}(t)]$ can be described by listing, for each p , the interval $[\underline{f}(p), \overline{f}(p)]$ of the possible quantile values:
 - the function $\underline{f}(p)$ is an inverse function to $\overline{F}(t)$,
and
 - the function $\overline{f}(p)$ is an inverse function to $\underline{F}(t)$.
- *Conclusion:* whatever method we have for DS knowledge bases, we can apply it to p-boxes as well.
- *Handling different types of uncertainty for different x_i :* just translate them into p-boxes.

Cauchy Method for Quadratic f

- *Linear case:* quadratic and higher order terms can be ignored.
- *Next case:* linear terms are still prevailing, but quadratic terms can no longer be ignored:

$$\delta y \stackrel{\text{def}}{=} \sum_{i=1}^n c_i \cdot \delta x_i + \sum_{i=1}^n \sum_{j=1}^n c_{ij} \cdot \delta x_i \cdot \delta x_j.$$

- *Analysis:* since linear terms are prevailing, max and min are attained when $\delta x_i = \pm \Delta_i$ (depending on $\varepsilon_i \stackrel{\text{def}}{=} \text{sign}(x_i)$):

$$\Delta^+ = \sum_{i=1}^n |c_i| \cdot \Delta_i + \sum_{i=1}^n \sum_{j=1}^n c_{ij} \cdot \varepsilon_i \cdot \varepsilon_j \cdot \Delta_i \cdot \Delta_j;$$

$$\Delta^- = \sum_{i=1}^n |c_i| \cdot \Delta_i - \sum_{i=1}^n \sum_{j=1}^n c_{ij} \cdot \varepsilon_i \cdot \varepsilon_j \cdot \Delta_i \cdot \Delta_j.$$

- *Problem:* for large n , literal computation takes too long.
- *Objective:* design a Cauchy-type method for this case.

Algorithm

- *Auxiliary algorithm:* $z = (z_1, \dots, z_n) \rightarrow g(z)$: apply the linear Cauchy deviate method to the auxiliary function $t \rightarrow \frac{1}{2} \cdot (f(x^{\text{mid}} + z + t) - f(x^{\text{mid}} + z - t))$ and the values $t_i \in [-\Delta_i, \Delta_i]$.
- *Main algorithm:*
 - We apply the algorithm $g(z)$ to the vector $0 = (0, \dots, 0)$, thus computing the value $g(0)$.
 - We apply the linear Cauchy deviate method to the auxiliary function

$$h(z) = \frac{1}{2} \cdot (g(z) - g(0) + f(x^{\text{mid}} + z) - f(x^{\text{mid}} - z));$$
 the result is the desired value Δ^+ .
 - Finally, we compute Δ^- as $2g(0) - \Delta^+$.
- *Computational complexity:* $2N^2$ calls to f .
- *Conclusion:* this method is better if $n \gg 8 \cdot 10^4$.

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